

Programming Quantum Computers: A Primer with IBM Q and D-Wave Exercises

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Quantum Programming Tutorial

Overview

- **Welcome**
- **Introduction to Quantum Computing (Patrick Dreher)**
 - Postulates of Quantum Mechanics, Linear Algebra, Qubits
 - Quirk Simulation
- **Gate-Level Quantum Computing (Greg Byrd)**
 - Quantum Gates, Circuits, and Algorithms
 - IBM Q Operation
 - IBM Q Programming with Qiskit
- **Adiabatic Quantum Computing (Frank Mueller)**
 - Basics of Quantum Annealing and QUBOs
 - D-Wave Programming
- **Programming Exercises with IBM Q and D-Wave**

What is a computer?

- Mathematical abstraction: a Turing machine
 - $M = \{Q, \Gamma, b, \Sigma, \delta, q_0, F\}$
 - All **states**, all symbols, blank symbol, input symbols, **transition function**, initial state, and final states
 - All of the preceding sets are finite, but the **memory** ("tape") on which they operate is **infinite**
 - Transition function
 - Maps {current state, symbol read} to {new state, symbol to write, left/right}
- Example: "If you're in state A and you see a 0, then write a 1, move to the left, and enter state B"



What else is a computer?

- **Nondeterministic Turing machine**
 - Replace transition function with a transition relation
 - Contradictions are allowed
 - Example: "If you're in state A and you see a 0, then simultaneously (i) write a 1, move to the left, and enter state B ; (ii) write a 0, move to the right, and enter state C ; and (iii) write a 1, move to the right, and enter state B ."
 - At each step, oracle suggests best path to take (unrealistic!)
- **Quantum Turing machine**
 - Same 7-tuple as in the base Turing machine
 - $M = \{Q, \Gamma, b, \Sigma, \delta, q_0, F\}$
 - **But...set of states is a Hilbert space**; alphabet is a (different) Hilbert space; blank symbol is a zero vector; transition function is a set of unitary matrices; initial state can be in a superposition of states; final state is a subspace of the Hilbert space
 - No change to input/output symbols; those stay classical

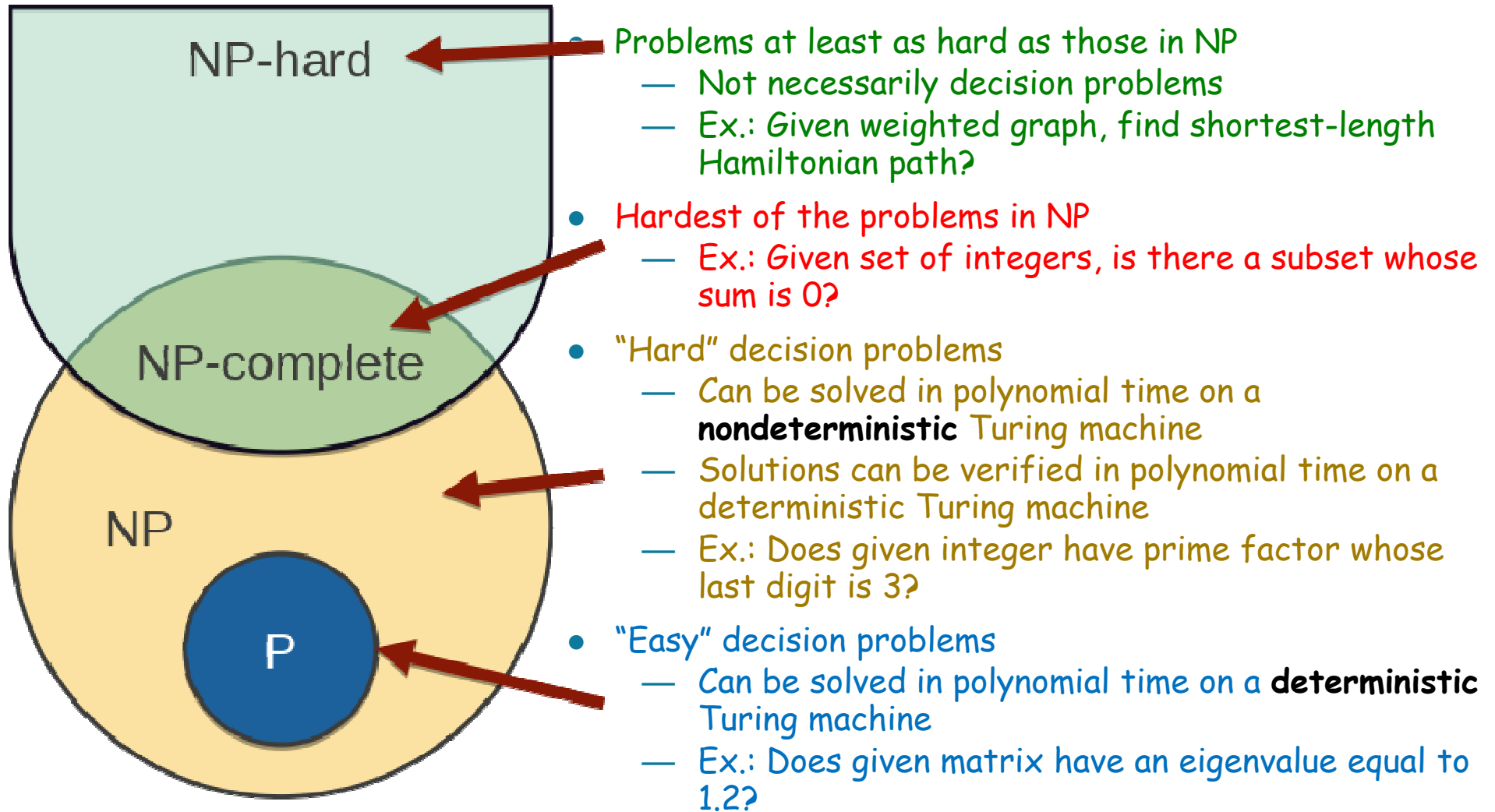
Introduction to Complexity Theory

- What problems can a computer solve quickly?
- Discuss in terms of **asymptotic complexity**, not wall-clock time
 - Ignore constants and all but the leading term
 - For input of size n , $O(n)$ can mean $3n$ seconds or $5n+2 \log n+3/n+20$ hours; it doesn't matter
 - **Polynomial time, $O(n^k)$ for any k , is considered good** (efficiently solvable), even if an input of size n takes $1000n^{20}$ years to complete
- **Superpolynomial** time—most commonly exponential time, $O(k^n)$ for $k>1$ —**is considered bad (intractable)**, even if an input of size n completes in only 2^n femtoseconds

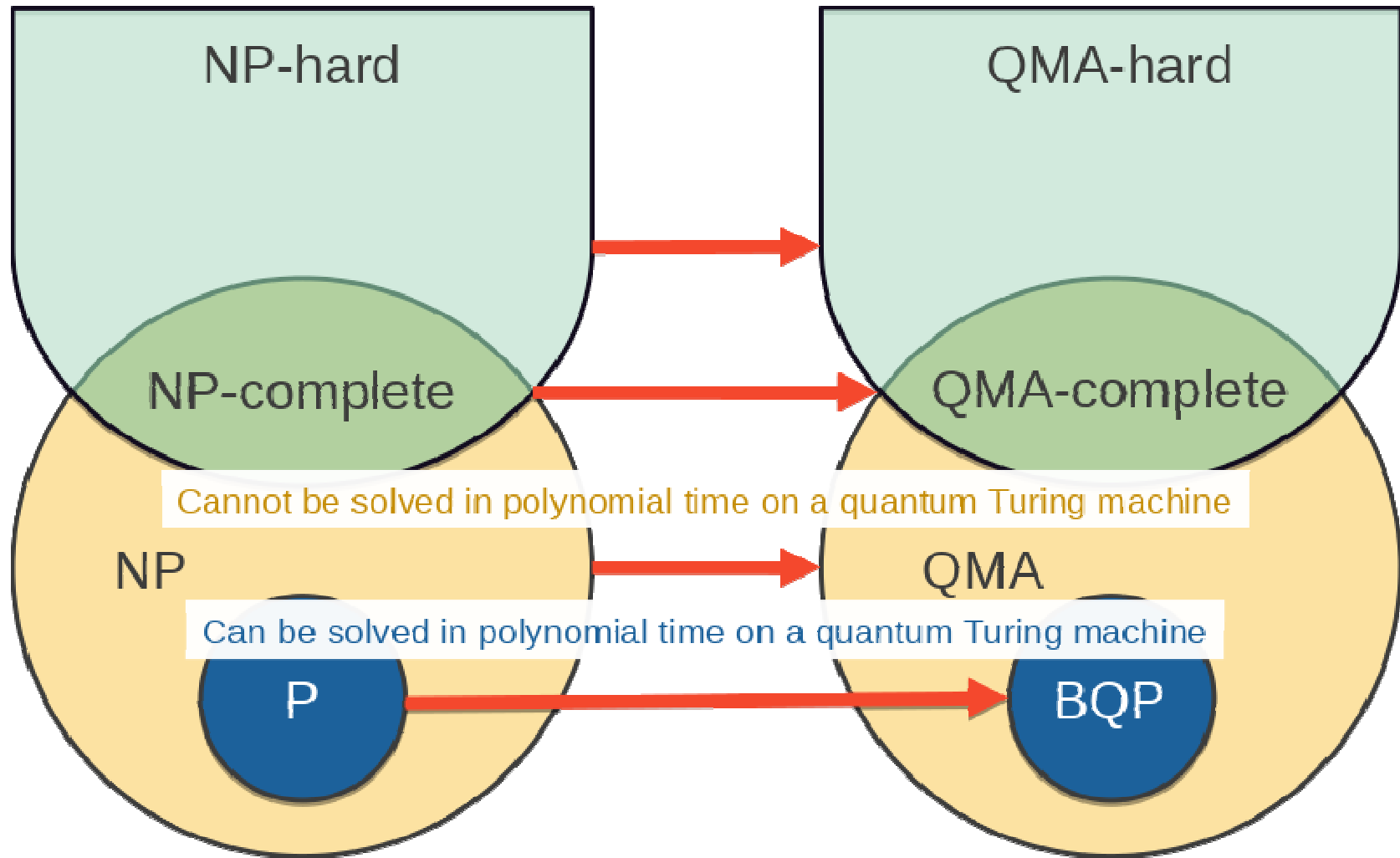
Introduction to Complexity Theory (cont.)

- Categorize problems into complexity classes
 - Goal: Determine which complexity classes are subsets or proper subsets of which other classes (i.e., representing, respectively, “no harder” or “easier” problems)
 - Approach is typically based on reductions: proofs that an efficient solution to a problem in one class implies an efficient solution to all problems in another class
- Typically focus on decision problems
 - Output is either “yes” or “no”

Venn Diagram of Common Complexity Classes

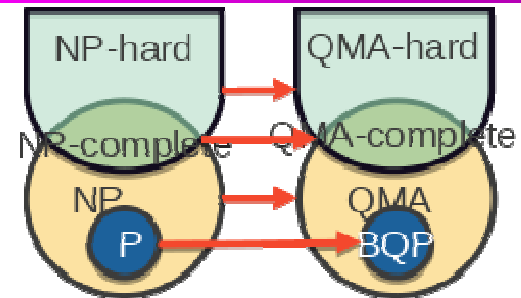


Quantum [Merlin Arthur] (QMA) Computing Complexity Classes



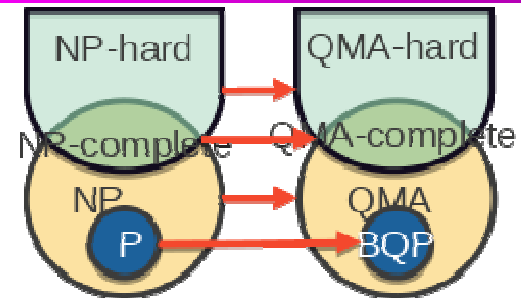
What Do We Know?

- Short answer: Almost nothing
- P vs. NP
 - $P \subseteq NP$
 - ??? $P = NP$ or $P \neq NP$; conjectured that $P \neq NP$
- NP-intermediate vs. NP-complete
 - NP-intermediate: set of problems in NP but not in NP-complete
 - $NP\text{-intermediate} \subseteq NP\text{-complete}$
 - ??? $NP\text{-intermediate} = NP\text{-complete}$
 - Implication: If $NP\text{-intermediate} \neq NP\text{-complete}$, then factoring (NP-intermediate) may in fact be an easy problem, but we just haven't found a good classical algorithm yet



What Do We Know? (cont.)

- P vs. BQP
 - $P \subseteq BQP$
 - ??? $P = BQP$ or $P \neq BQP$
 - Implication: If $P = BQP$, then quantum offer no substantial (i.e., superpolynomial) performance advantage over classical
- NP-complete vs. BQP
 - ??? BQP vs. NP-complete; conjectured $BQP \subset NP\text{-complete}$
 - Implication: Believed that quantum computers cannot solve NP-complete problems in polynomial time
- Initial focus: Quantum supremacy \rightarrow break complexity class
- Today's focus: Quantum advantage \rightarrow faster than classical
 - By constant factor



It's Not All Doom and Gloom

- Sure, quantum computers probably can't solve NP-complete problems in polynomial time
- Still, even a polynomial-time improvement is better than nothing
- Grover's algorithm
 - Find an item in an unordered list
 - $O(n) \rightarrow O(\sqrt{n})$
- Shor's algorithm
 - Factor an integer into primes (NP, but not NP-complete)
 - $O(2^{\sqrt[3]{n}}) \rightarrow O((\log n)^3)$

Quantum Architectures

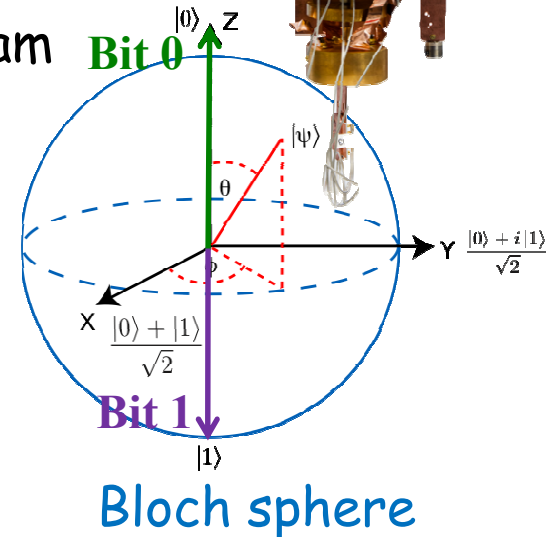
1. Quantum annealer (D-Wave)
 - Specialized: optimization problems → find lowest energy level
 - Uses tunneling and entanglement
 - Better than classical? → unknown, maybe significant speedup
2. Approximate quantum [gate] computer (IBM Q, Regetti, IonQ...)
 - More general: optimization, quantum chemistry, machine learning
 - Superposition, entanglement
 - Better than classical? → likely, sign. Speedup: "advantage"
3. Fault-tolerant quantum computer (in "some years" from now)
 - Deals w/ errors (noise) algorithmically
 - Most general: crypto, search, and any of the above ones
 - Need 1000 physical qubits per virtual ("error-free") qubit
 - Better than classical? → proved theoretically: "supremacy"

Quantum Algorithms (Gate Model)

- Key concepts
 - N classical bits go in, N classical bits come out
 - Can **operate on all 2^N possibilities** in between
 - Requirement: Computation must be reversible (not a big deal in practice)
 - Main challenge: You get only **one measurement**; how do you know to measure the answer you're looking for?
 - High-level approach: Quantum states based on complex-valued probability amplitudes, not probabilities—can sum to 0 to make a possibility go away
 - Very difficult in practice
- Google "quantum algorithm zoo" → 60 algorithms known to date
 - Based on only a handful of building blocks
 - Each requires substantial cleverness; not much in the way of a standard approach

Gate model (cont.)

- Examples: IBM Q, Regatti, IonQ, Intel, Google...
- Programming = set parameters of physics experiment, use lasers/radio freq. to energize qubits, observe results
 - Lasers/radio freq. triggered by your program
 - Program = circuit of basic quantum gates
 - Quantum: CNOT ..., classical: NAND ...
 - Clock rate in us range
- 2^N states \rightarrow qubits in "superposition"
 - IBM Q: 20 qubits $\rightarrow 2^{20}$ states today
 - Qubit: $|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$, $|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ as column vector \rightarrow
 - Superposition: 0 & 1 "at the same time" $|\psi\rangle = a|0\rangle + b|1\rangle$, $|a|^2 + |b|^2 = 1$
 - Example: 3 qubits, overall state $|\psi\rangle = a|000\rangle + b|001\rangle + c|010\rangle \dots$
 - Repeat measurement \rightarrow probability per state: $|a|^2, |b|^2, |c|^2$
 - new results every few ms



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