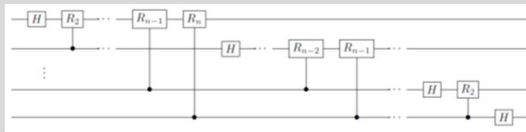


# Quantum Fourier Transform

ECE 592/CSC 591 – Fall 2018



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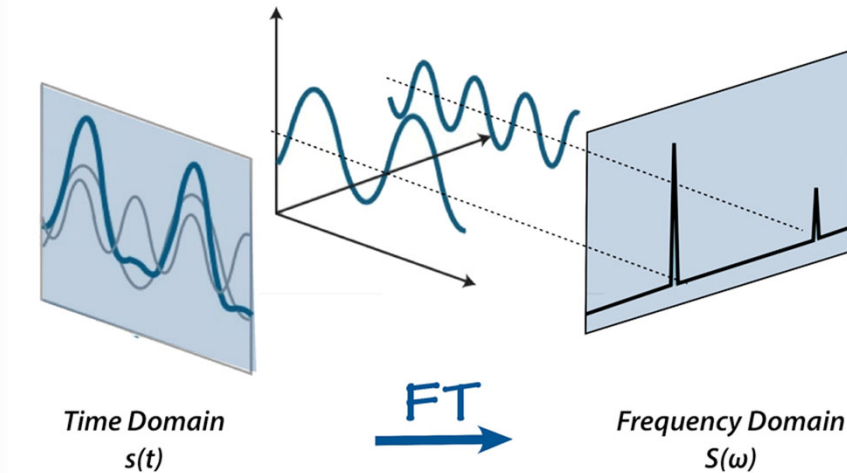
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## QFT

- Quantum version of FFT
- Building block of other quantum algorithms
- Transforms coefficients of  $n$ -qubit superposition into “frequency domain” – useful in period finding
- $N$ -dimensional ( $n$ -qubit) QFT:  $O(\log^2 N)$ 
  - Classical FFT:  $O(N \log N)$
  - Exponential speedup

## Review: Fourier Transform



<http://mriquestions.com/fourier-transform-ft.html>

## DFT and FFT

DFT operates on a discrete complex-valued function  $a(x)$ , and produces a new function  $A(x)$ :

$$A(x) = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} a(k) e^{2\pi i \frac{kx}{N}}$$

FFT is an efficient implementation of DFT when  $N$  is a power of 2.

## FFT

Let  $\omega_{(n)}$  represent the  $N$ th root of unity,  $\omega_{(n)} = e^{\frac{2\pi i}{N}}$ .

FFT can be expressed as a matrix transformation

$$(F_N)_{jk} = \omega_{(n)}^{jk}$$

Example, for  $N = 4$ ,  $F_4$  is a  $4 \times 4$  matrix:

$$F_4 = \frac{1}{\sqrt{4}} \begin{pmatrix} \omega_{(2)}^0 & \omega_{(2)}^0 & \omega_{(2)}^0 & \omega_{(2)}^0 \\ \omega_{(2)}^0 & \omega_{(2)}^1 & \omega_{(2)}^2 & \omega_{(2)}^3 \\ \omega_{(2)}^0 & \omega_{(2)}^2 & \omega_{(2)}^4 & \omega_{(2)}^6 \\ \omega_{(2)}^0 & \omega_{(2)}^3 & \omega_{(2)}^6 & \omega_{(2)}^9 \end{pmatrix} = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix}$$

Implementation efficiency comes from recursive nature of FFT – more about this later.

## QFT uses the same matrix

$$\sum_x a(x) |x\rangle \rightarrow \sum_x A(x) |x\rangle$$

Function is expressed as coefficient of  $n$ -qubit state  
No output register – change in state.

Same matrix, normalization makes it unitary

$$\text{QFT}_4 = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix} \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{pmatrix} \begin{pmatrix} a(0) \\ a(1) \\ a(2) \\ a(3) \end{pmatrix} = \begin{pmatrix} A(0) \\ A(1) \\ A(2) \\ A(3) \end{pmatrix}$$

## Effects of QFT

If  $a_x$  is periodic with period  $r$ , which is a power of 2, then  $A_x$  will be zero except when  $x$  is a multiple of  $\frac{N}{r}$ .

$$\sum_x a(x) |x\rangle \rightarrow \sum_x A(x) |x\rangle$$

When measured, resulting state will be a random multiple of  $\frac{N}{r}$ .

If period is not a power of 2, then states near multiples of  $\frac{N}{r}$  will be measured with high probability.

## With one qubit...

- One-qubit QFT = Hadamard gate.

$$\omega_{(1)} = e^{\frac{2\pi i}{2}} = -1$$

$$\text{QFT}_2 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

# Implementation

- To the "chalkboard..."

# Effect

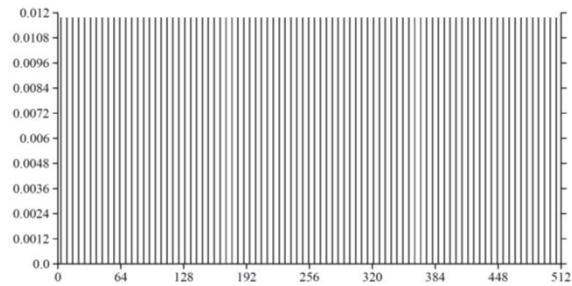


Fig. 2. Probabilities for measuring  $x$  when measuring the state  $C \sum_{x \in X} |x, 8\rangle$  obtained in Step 2, where  $X = \{x | 211^x \bmod 21 = 8\}$

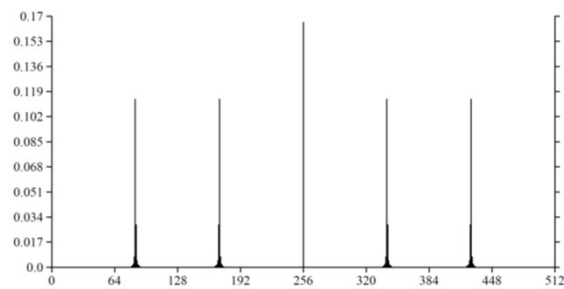


Fig. 3. Probability distribution of the quantum state after Fourier Transformation.