

# **Driving to the 48 USA State Capitals: Programming the D-Wave QPU**



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# Overview

- *Review* (Scott Pakin and D-Wave)
- qbsolv
- 7-city Traveling Salesman Problem
- 48-city Traveling Salesman Problem
- Summary

# Review: Programming Model

QUBIT		Quantum bit which participates in annealing cycle and settles into one of two possible final states:
COUPLER		Physical device that allows one <b>qubit</b> to influence another <b>qubit</b>
WEIGHT		Real-valued constant associated with each <b>qubit</b> , which influences the <b>qubit's</b> tendency to collapse into its two possible final states; controlled by the programmer
STRENGTH		Real-valued constant associated with each <b>coupler</b> , which controls the influence exerted by one <b>qubit</b> on another; controlled by the programmer
OBJECTIVE		Real-valued function which is minimized during the annealing cycle

$$Obj(a_i, b_{ij}; q_i) = \sum_i a_i q_i + \sum_{ij} b_{ij} q_i q_j$$

The system **samples** from the  $q_i$  that minimize the objective



# Review: The overall process

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- Need to map problems into binary variables
- Need to map the binary variable expressions into linear terms and pair quadratic terms
  - QUBO: Quadratic Unconstrained Binary Optimization
- Run the QUBO on the D-Wave QPU
- Interpret the results
- Revisit the mapping...

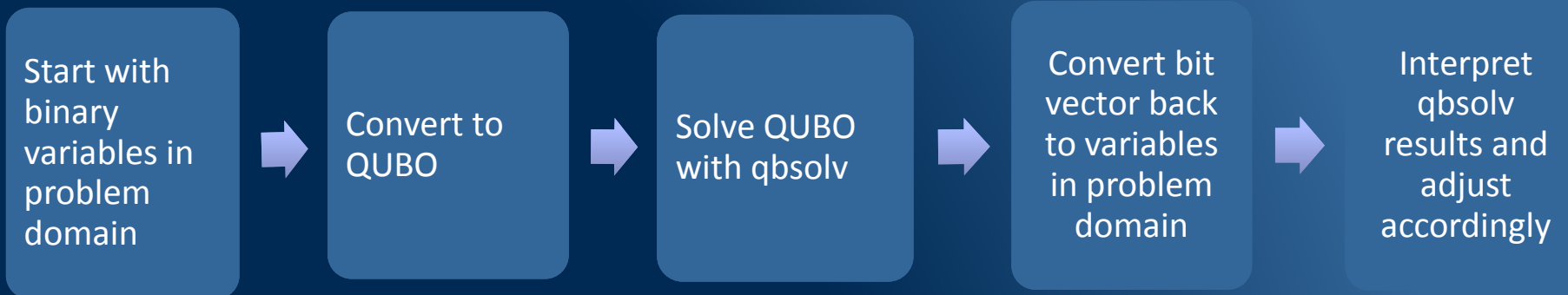


# QUBOs can be found in many fields

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- Optimization
  - Recent: financial portfolio management
  - Recent: hospital re-admission statistics (Medicare funding)
  - Recent: bioinformatics / Multiple Sequence Alignment
  - Recent: nuclear power plant failure scenario analysis
  - Mathematical problems like Minimum Vertex Cover
  - Job-shop scheduling, other graph problems (Map Coloring, vertex set color)
- Sampling (from probability distribution)

# Application example flow



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# The need for qbsolv

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- Many problems require many more qubits or couplers than are available with the current chip.
- Examples:
  - Portfolio Optimization example requires 63-variable complete graph, but 2000Q does not go beyond ~50-qubit complete graph with direct embedding
  - 48-city Traveling Salesman requires complete graph, and it might be possible on the 2000Q, but the needed embedding would have very long chains, which do not perform well
  - Traffic flow optimization example using 418 cars (with 3 routes each) require almost 2000 highly-connected variables, too many to fit directly on the 2000Q chip.





# qbsolv

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- Hybrid quantum/classical QUBO solver (tabu = classical heuristic solver)
- Designed for problems too large and/or too dense to run on D-Wave quantum computer
- Divides problems into chunks, and iterates on sub-QUBOs (similar to HFS algorithm)
- Open-source: <https://github.com/dwavesystems/qbsolv>
- Can be used standalone, or with 128-qubit simulator, or with QPU
- Produces a single bitstring solution representing the final states of all the binary variables

# Motivating Algorithm

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- "A Multilevel Algorithm for Large Unconstrained Binary Quadratic Optimization", Wang, Lu, Glover, and Hao [2012]
- Principles
  - Identify the *backbone* of the QUBO; *i.e.*, the variable settings that are correlated for all valid answers, or, contribute the most to a local optimum
  - Select subQUBOs by most/least impact of each variable in determining the optimum
  - Solve the subQUBOs with a solver known to run effectively at smaller scale
  - Propagate subQUBO answer out to original variables in full QUBO
  - Iterate above steps until no further improvement



# How qbsolv works

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- Hybrid algorithm:
  - Identify the significant rows and columns of the larger problem (What changes a lot with spin flips? What doesn't?)
  - Create a smaller representative QUBO of that subset
  - Execute that smaller QUBO on the D-Wave system (pre-computed embedding, speeds up run-time)
  - Use the answer to guide the larger solver (new starting point, closer to the minimum)

# Example qbsolv Output

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```
$ qbsolv -i bqp50.qubo1.qubo
```

```
50 Number of bits in solution
```

```
10111111101001111101001101011111101111111110110110
```

```
-5176.00000 Energy of solution
```

```
0 Number of Partitioned calls
```

```
0.21352 seconds of classic cpu time
```

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# 7-city Traveling Salesman Problem

- A learning exercise
- We will explore:
  1. QUBO for leg variables
  2. City visit constraints
  3. Embedding and chain constraints
  4. Importance of parameters and problem construction



# Problem Specifications

## Optimization:

- Given a list of cities and the distances between each pair of cities, what is the shortest possible route that visits each city exactly once and returns to the origin city?
- Given a length  $L$ , decide whether the list of cities and distances has any tour shorter than  $L$ .
- Symmetric TSP; Undirected graph, distances in miles; distances obtained from various Web sites .

# The Cities

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A = Albuquerque, NM

B = Boston, MA

C = Charlotte, NC

D = Detroit, MI

E = Evanston, IL

F = Frankfort, KY

G = Gulfport, MS

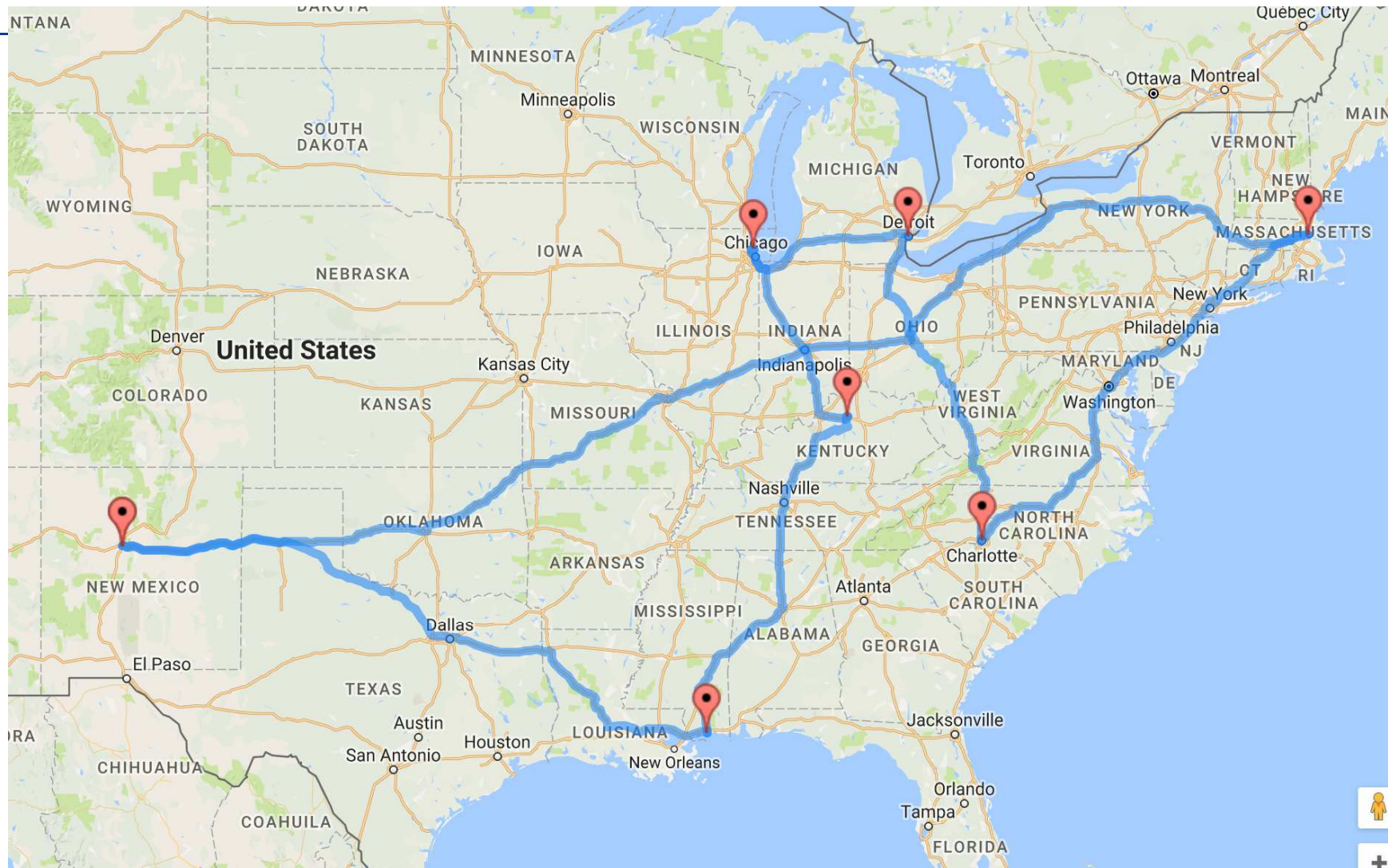
(Distances found in file tsp7.b)

Example path:

Path = (A->B) + (B->C) + (C->D) + (D->E) + (E->F) + (F->G) + (G->A)



# Map of Alphabetical Order Path



# QUBO: Leg Variables

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Binary Variable  $ab$ :

1 if the trip includes the segment A  $\rightarrow$  B

0 if the trip does not include A  $\rightarrow$  B

Distances between cities A and B denoted by  $D_{ab}$

Distance to Minimize:

$$D_{ab} * ab + D_{ac} * ac + D_{ad} * ad + \dots + f_g * D_{fg}$$

How do we convert this into a QUBO?

# Converting into a QUBO

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Each city must be visited exactly twice – once arriving, and once departing.

For city A, we must have:

$$ab + ac + ad + ae + af + ag - 2 = 0$$

For City B,

$$ab + bc + bd + be + bf + bg - 2 = 0$$

And so on, up to city G.

The QUBO to minimize:

$$D_{ab} * ab + D_{ac} * ac + D_{ad} * ad + \dots + D_{fg} * fg + \gamma((ab+ac+ad+\dots-2)**2 + (ab + bc + bd + be + bf + bg - 2)**2 + \dots)$$

# Algebraic results: expanding the equation

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- Groups of terms of the form:
  - $3 * vertex * ab$ : favor visiting the path A->B
- Groups of terms of the form:
  - $2 * vertex * ab * ag$ : penalize visiting the paths A->B and A->G  
(some of these will be selected)
- There are seven explicit **dw** assert statements (which help identify valid solutions):
  - $assert:vertex:ab + ac + ad + ae + af + ag - 2$
- *vertex* is a Lagrangian multiplier; needs to be weighed against the inter-city distances

# Steps to run the 7-city TSP (write a run script)

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- Prepared `tsp7.b` (parameter file) and `tsp7.q` (QUBO file) by hand
- Embed the problem onto the 128-qubit simulator using the **dw embed** command (the default embedder algorithm tries to find a way to map the needed logical qubits onto available physical qubits)
- Try a value of `vertex` and a value of `param_chain`, another adjustable parameter (controls the chain strength in the embedding)
- Run **dw bind** to bind the parameter values in the B file to the QUBO
- Run **dw exec** to run the problem on the simulator (or QPU)
- Run **dw val** to interpret and validate the output

# Parameter explorations

- Earlier, we mentioned *vertex* and *param\_chain*, the adjustable parameters
- *Vertex* controls the strength of obeying the constraints
- *Param\_chain* controls the strength of chains in the embedding
- When running the problem, parameter space must be explored, to find largest possible number of solutions

	1300	1400	1500
4500			
4750			
5000			

# Looking at the lowest-energy solution

- **dw val -s 1 tsp7.sol** produces output that looks like this:

```
**** SOLUTION 1 ****
```

```
ab <== 0
```

```
af <== 1
```

```
ag <== 1
```

```
bc <== 1
```

```
...
```

```
VALID:    Y
```

```
SAMPLES:  4
```

```
OBJECTIVE: -33778.00
```

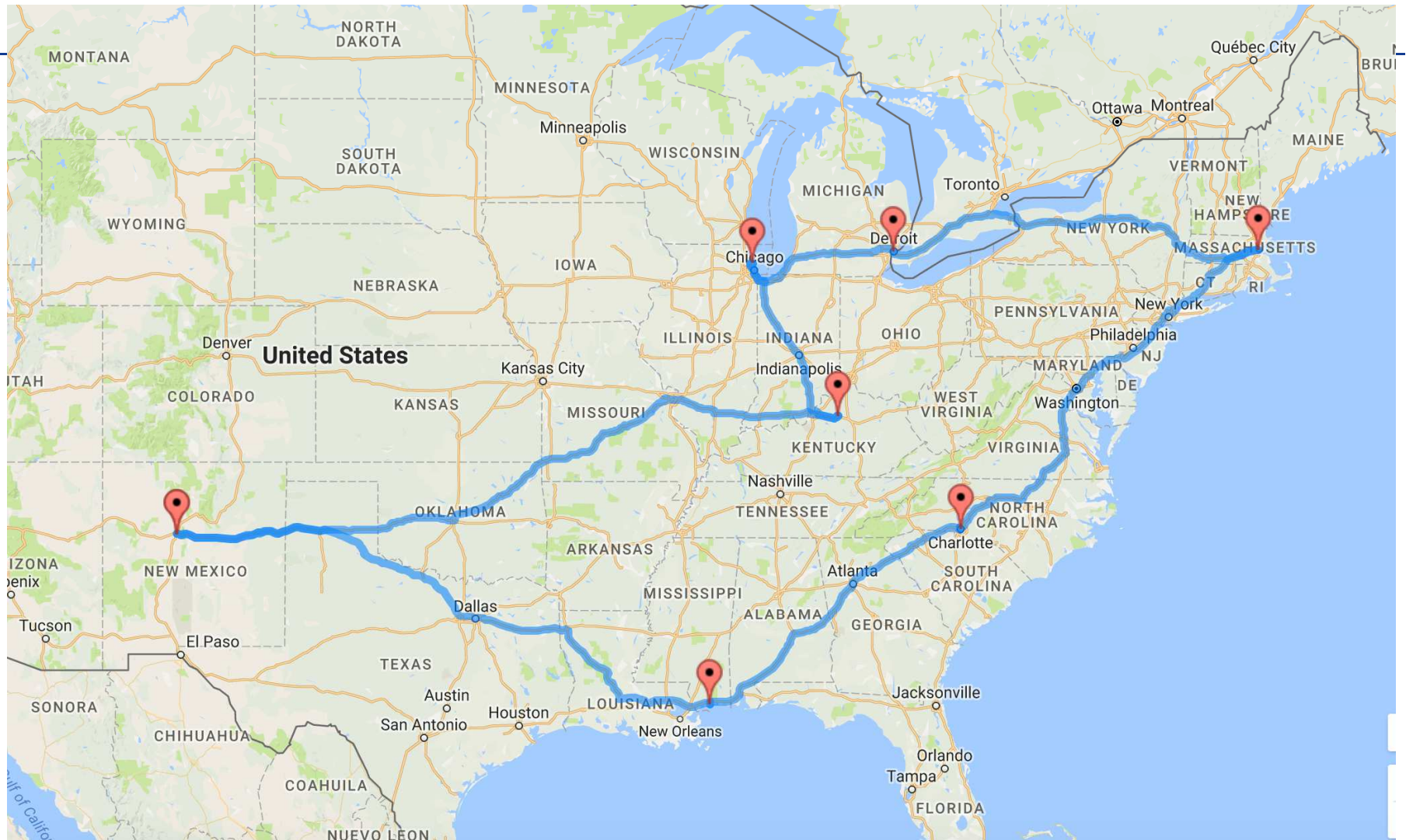
And we get the path:

*af + ef + de + bd + bc + cg + ag*

*a -> f -> e -> d -> b -> c -> g -> a*



# A map of the lowest-energy solution





# Interpreting the energy

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The path:  $af + ef + de + bd + bc + cg + ag$

Adding up the miles: 5422 miles

**dw** tells us the energy is -33778

For each vertex, there will be two nonzero terms of the form:

$$-3 * vertex * af = -3 * vertex$$

And one term of the form:

$$2 * vertex * af * ag = 2 * vertex$$

Thus:  $-4 * vertex$  per 7 cities  $\Rightarrow -28 * vertex$

$$Energy = Total Mileage - 28 * vertex$$

(for vertex = 1400, this equality works)

# An interesting issue

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A valid solution and different path:

*ae ag bc bd cf df eg*

Notice:  $A \rightarrow E \rightarrow G \rightarrow A$

This is called a **subloop**: the path split into a 3-loop and a 4-loop

- Subloops are not prevented by “visit twice” constraint
- Subloops are not desirable solutions; we would have to write many more constraints to eliminate them
- The Lucas formulation (**Permutation matrix**) rules out subloops (next section)



## Conclusions from 7-city TSP

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- **dw** can be used as part of overall toolbox, start to finish
- Good problem for exploring Lagrangian parameters, understanding solutions, how QUBO construction affects solutions (subloop problem).
- Good problem for understanding chains and embedding
- Has quick physical interpretation of solution
- Good problem for understanding what the QPU does, what its inputs are, and what it returns

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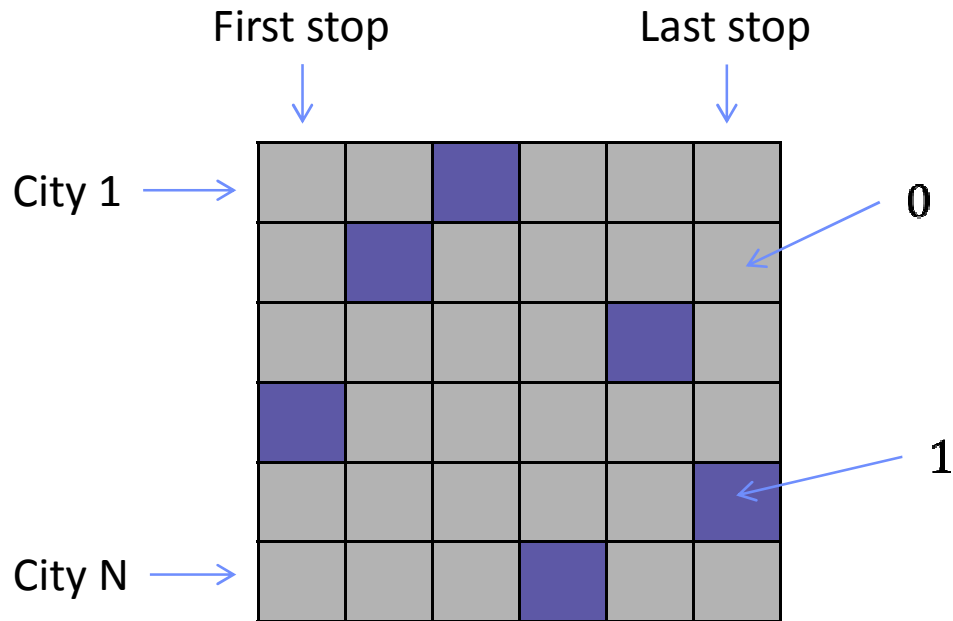
# 48-city Problem Specifications

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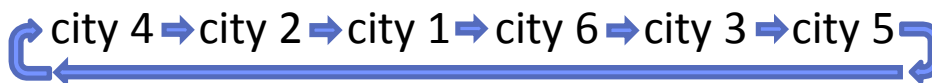
- Symmetric TSP; undirected graph; distances in miles
- Driving distances obtained from Google Maps
- We can reduce the complexity of the problem by 1- assume last city visited is alphabetically last city (Cheyenne, WY)
- Andrew Lucas paper: <https://arxiv.org/pdf/1302.5843.pdf>

# First approach: QUBO (ref. Andrew Lucas)

QUBO approach means building a **permutation matrix** of 0's and 1's and then introducing quadratic terms to include distances:



Tour represented:



## QUBO (continued)

For each row and column, we introduce a constraint via the following QUBO terms:

$$\text{Constraint} = A \left( -1 + \sum_{i=1}^N x_{1,i} \right)^2 \quad \leftarrow \text{For row 1}$$

$$\text{Constraint} = A \left( -1 + \sum_{j=1}^N x_{j,1} \right)^2 \quad \leftarrow \text{For column 1}$$

Distance from the first city visited to the second city visited is computed like this:

$$\text{Distance} = B \sum_{i=1}^N \sum_{j=1}^N d_{i,j} x_{i,1} x_{j,2} \quad \leftarrow \text{From 1}^{\text{st}} \text{ to } 2^{\text{nd}} \text{ city}$$

# Algebraic results

- QUBO approach requires boolean variables to encode a tour, but “last city visited” assumption provides reduction
- Constant term  $(2N-2)A$ : one for each row + column, reduced
- Each variable gets a term  $-2A$  multiplied by it (e.g.  $-2A$ ) to incentivize visiting it on a particular step (e.g., city 1, step  $i$ )
- Pairwise terms penalize cities being visited twice, or visited on same step (e.g.  $2A$  or  $2A$ )
- Distance terms penalize following an edge between two cities (the distance term raises the overall energy) (e.g.  $B$ )
- “Last city visited” assumption leads to diagonal terms  $B$  and  $B$  (last and first have to get to Cheyenne)



# Python program generate\_qubo.py

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- Number of cities  $N$  is input, cannot go beyond 48, but easy to extend
- Divided through by  $B$ , so that there is only one adjustable parameter  $A$
- Read in the inter-city distances from state\_capitals.txt, and create distance matrix
- Diagonal terms will be  $-2A$ , except when we need to include “last city visited” reduction effect
- qbsolv requires off-diagonal terms to have  $i < j$
- Use QUBO\_details library to write QUBO



# Looking at the 48-city QUBO

- For  $N = 48$ , with the reduction,  $(N-1)^2$  Boolean variables (2209)
- A set to 8500 (required: larger than biggest inter-city distance)
- Most of the diagonals are  $-17000$  ( $-2 * A$ ), but some have the distance added in, to the last city (“How far from here to Cheyenne?”)
- Many of the off-diagonal terms are  $17000$  ( $2 * A$ ) as well
- The last off-diagonal term is (2207,2208), as we expect

# Python program interpret\_lucas.py

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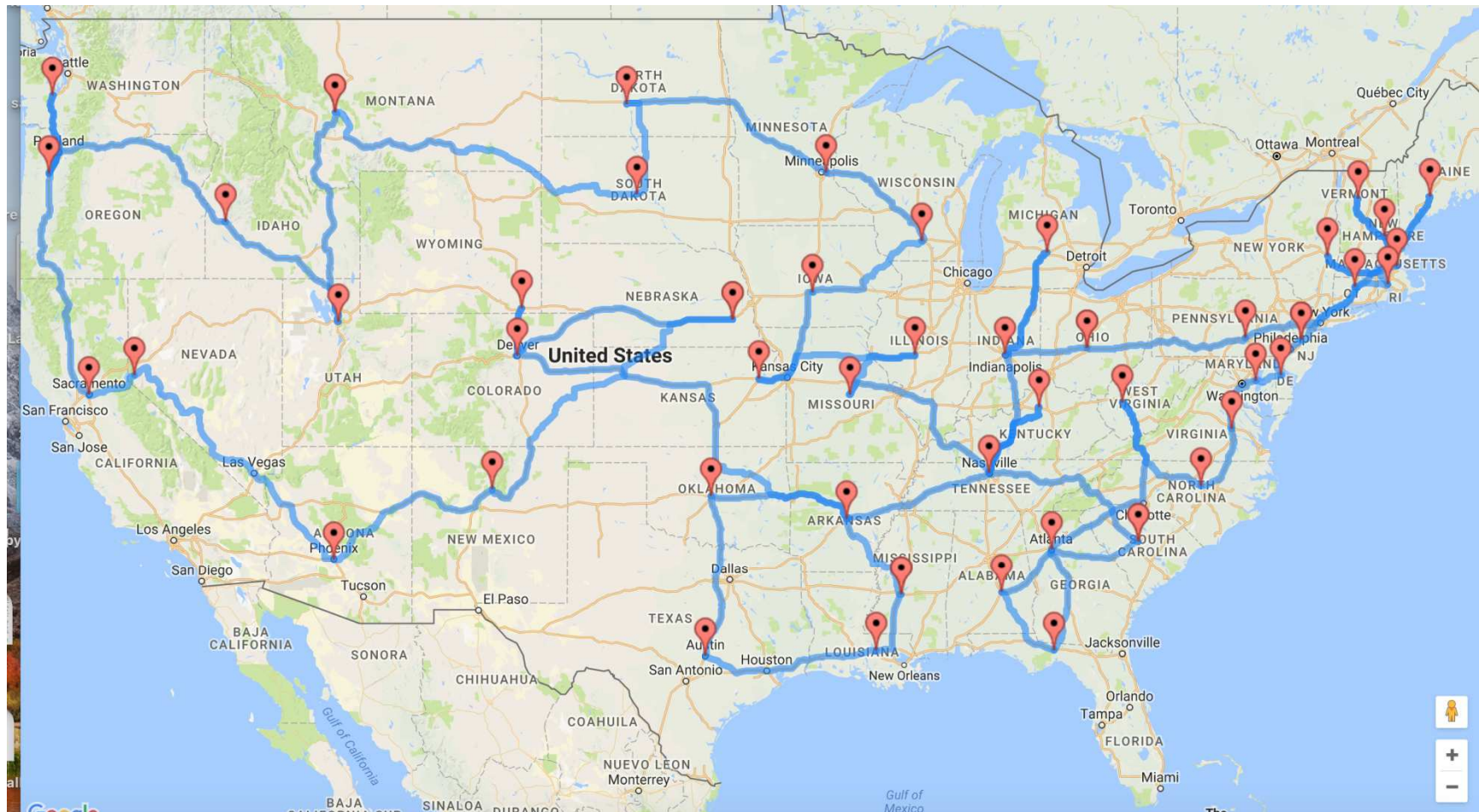
- Read in the inter-city distances from state\_capitals.txt, and create distance matrix
- Read in the names of the cities
- Read in the qbsolv output
- Associate the result bitstring with binary variables to compute the city tour
- Output a city tour for eventual drawing on Google Maps
- Useful for debugging QUBO/energy

# Python program generate\_map.py

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- Read in the city tour from the previous step
- Read in an HTML template for generating Google USA map
- Insert the city tour into the HTML template
- Write out an HTML page which can be displayed, containing the Google map of the route!

# The solution – meh



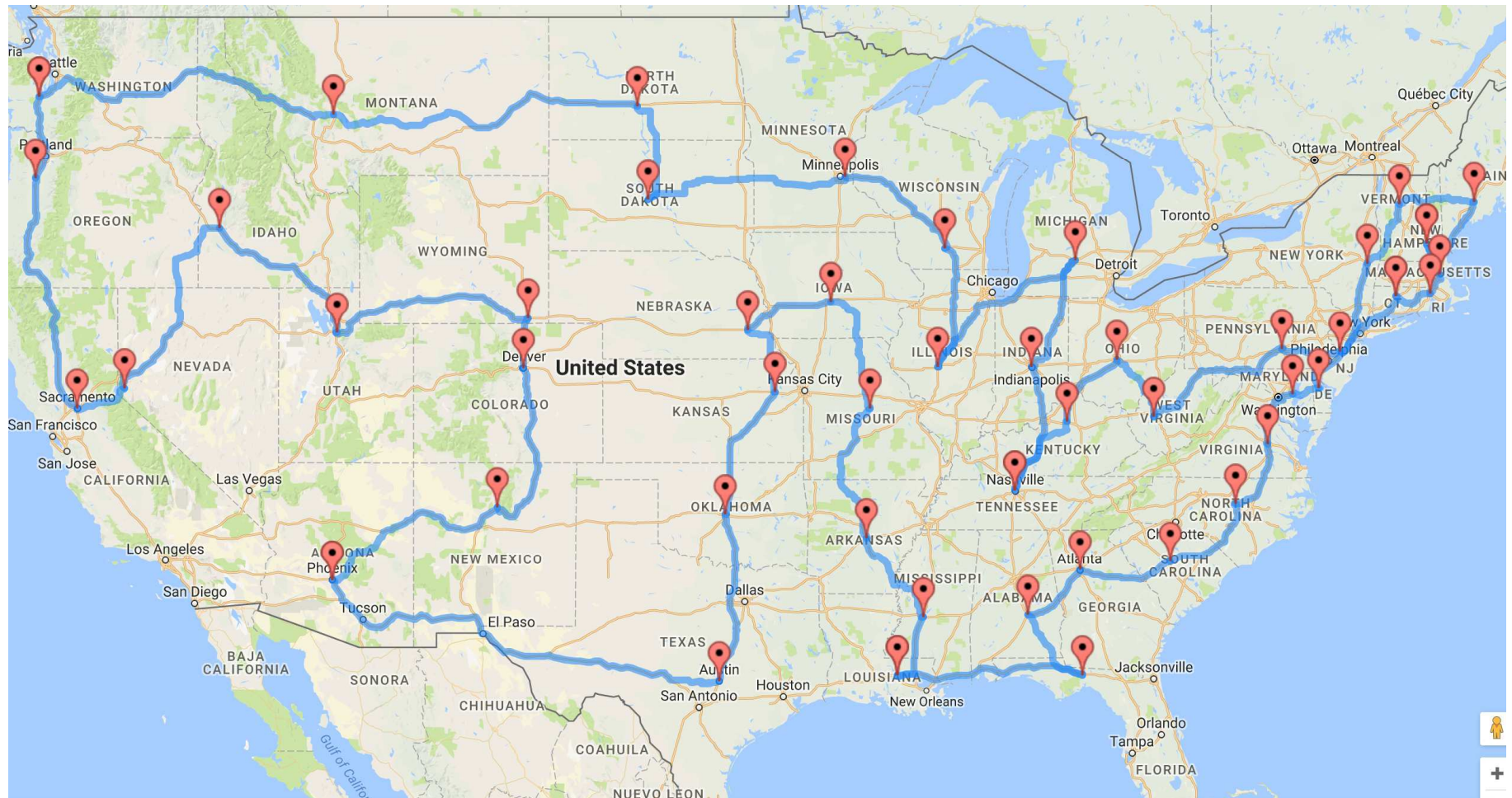


# A possible “Optimal solution” from Randal Olson

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- Randal Olson from the University of Pennsylvania wanted to drive all 48 state capitol buildings
- Genetic algorithm
- Olson’s Web page indicates his route is 13,310 miles, but he focuses on state capitol buildings
- <http://www.randalolson.com/2016/06/05/computing-optimal-road-trips-on-a-limited-budget/>

# Randal Olson's solution – way better





# Insights

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- Why didn't we get the better solution?
- Notice that the difference between "good" and "bad" solutions is less than 1% of the overall energy – why is this?
- Some problems are mathematically not great for the D-Wave
- Only  $N-1=47$  1's can be turned on; the qbsolv algorithms are more effective with larger numbers of bitflips (I need to try this the other way)
- For your problems, focus initially on special cases which can be intuitively understood (TSP with 4 cities, etc.)
- Use post-qbsolv code to test solutions for consistency (for example, introduce additional parameters and establish that the solutions don't change)



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# Summary

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- Hybrid classical/quantum approach to using D-Wave
- “Toolbox” approach: Python, dw, qbsolv, simulator, (hardware)
- Problem has understandable graphical representation
- “This problem is not particularly large, but hard.”
- The goal: you can formulate your client’s problems into QUBO and then run them, even if solutions are not immediately optimal
- The post-qbsolv code is vital for at least three reasons:
  - Interpreting the bitstrings and understanding the solutions
  - Debugging the math of converting the problem to QUBO
  - Debugging the code representing the QUBO