



# CoMerge

Toward Efficient Data Placement in Shared Heterogeneous

Memory Systems

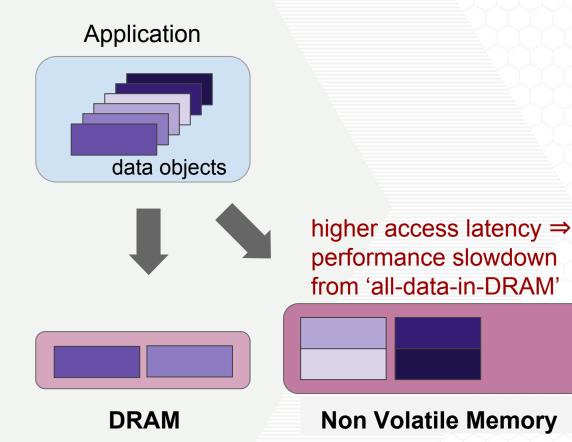
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### **Motivation**

Performance slowdown in heterogeneous memory systems.





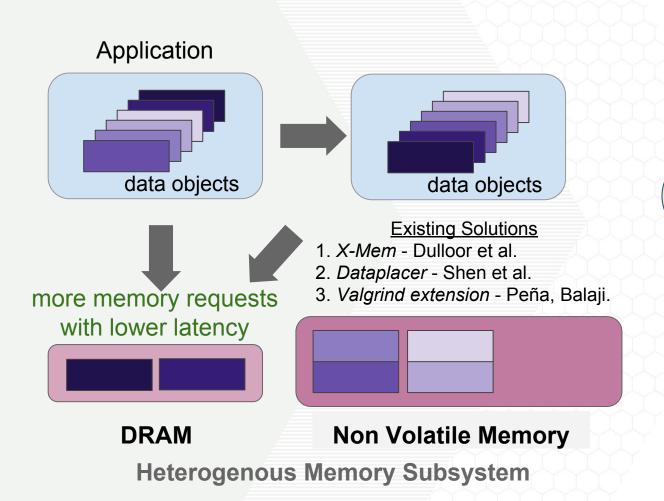
**Heterogenous Memory Subsystem** 



### **Existing Solutions**







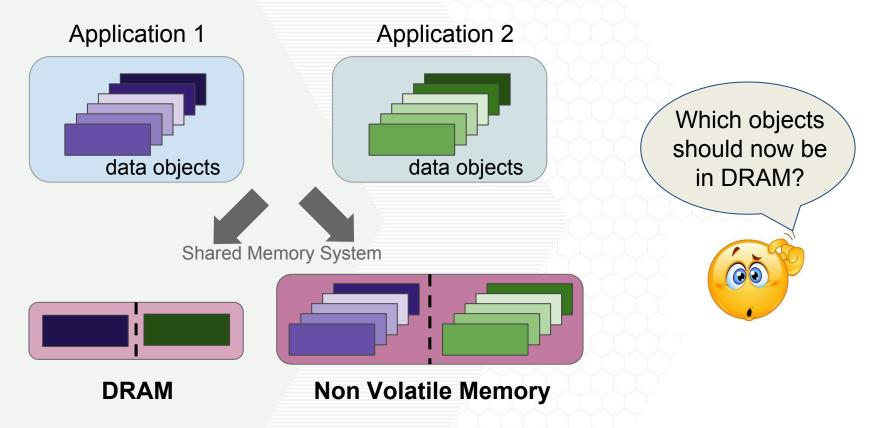
Think about which objects get allocated in DRAM.



### **Problem Statement**

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Limited Utility of Existing Solutions in Shared Systems.



Do the partitioning techniques using existing solutions:

- Reduce the slowdown across all collocated applications?
- Maximize DRAM utilization?

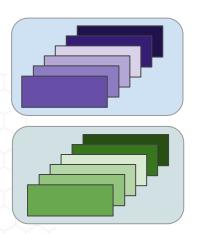


### **Our Contributions**

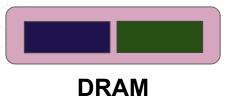
What do we need to do differently?



- Sorting objects within one application:
   co-benefit metric captures:
  - Exact contribution of a data object to overall application runtime.
  - Overall application sensitivity to execution over Non-Volatile Memory.



- 2. Distributing DRAM across applications: **CoMerge** memory sharing technique.
  - a. Mitigates slowdown across all collocated applications.
  - b. Maximizes the DRAM usage.



### **Observations**

What are we going to see next?



- 1. Not all applications are slowed down in the same degree when accessing Non Volatile Memory.
- 2. Not all data objects of an application help reduce the performance slowdown, when placed in DRAM.

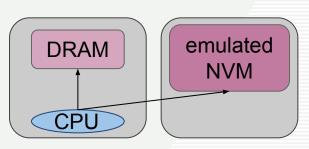
### **Polybench** Benchmarks

- 30 simple algebraic kernels.
- Single-threaded.

### **CORAL** Suite of mini-apps

- 3 HPC representative kernels.
- Multi-threaded. OpenMP.

#### Hardware Testbed



Emulate Non Volatile Memory for various combinations of *reduced bandwidth* and *increased latency*.

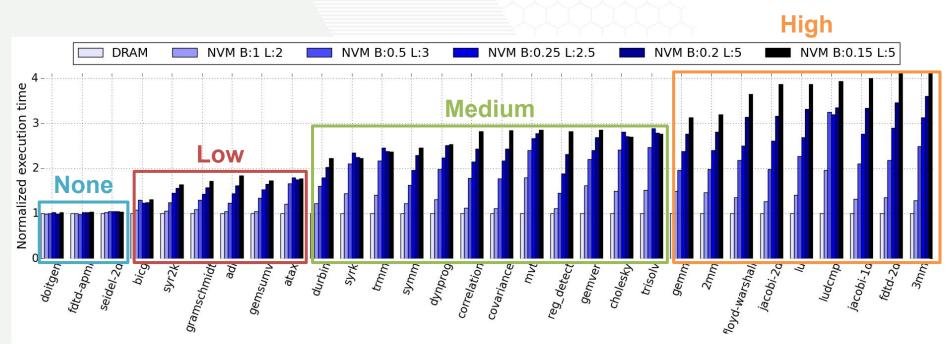
e.g. B 0.5 : L 2

0.5 times less bandwidth: 2 times more latency

# **Overall Application Sensitivity**



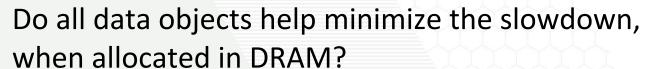
Do all applications get slowed down in the same way when accessing Non Volatile Memory?



Performance slowdown across Polybench/C, normalized to 'all-data-in-DRAM' execution.

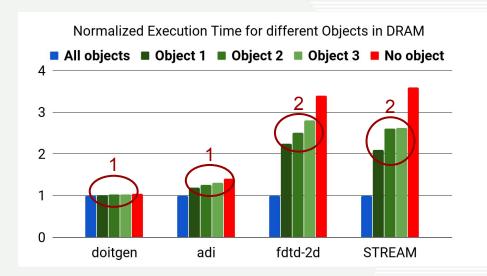
Applications show different levels of sensitivity to execution over slower memory components.

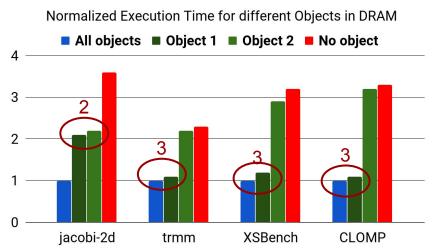
# **Data Object Sensitivity**





fixed NVM at B 0.2: L 5





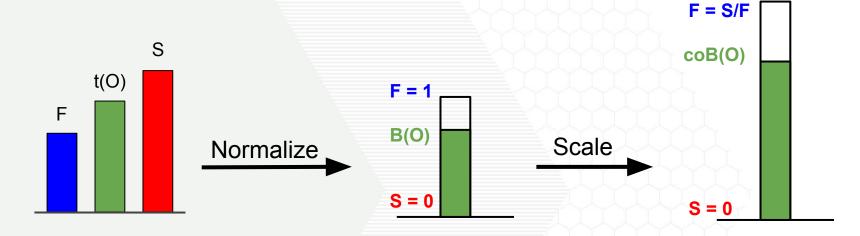
#### <u>Observations</u>

- 1. For non or low sensitive apps, doesn't matter which object is in DRAM.
- 2. Different data objects can contribute equally to the application runtime.
- 3. There can be objects whose allocation in DRAM is the only way to mitigate slowdown.

### **Co-Benefit Metric**

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Can we capture the previous observations?



Run Time	Objects in DRAM
F	All
t(O)	object O
S	None

How much does a specific object help reduce the slowdown?

How can we make sure that objects of higher sensitivity kernels are getting prioritized?

e.g. 
$$B(O) = 0.9$$

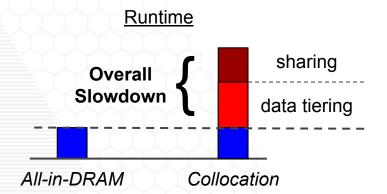
$$coB(O) = 0.9 * low sensitivity = 0.9$$
  
 $coB(O) = 0.9 * high sensitivity = 3.9$ 

### **DRAM Distribution**

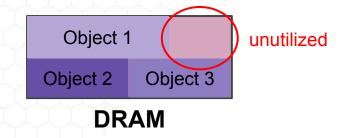


What are the goals of an efficient technique?

1. Minimize overall runtime slowdown across all applications.



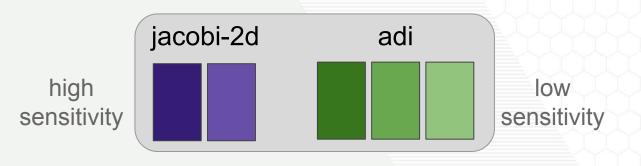
2. Maximize the utilization of DRAM.



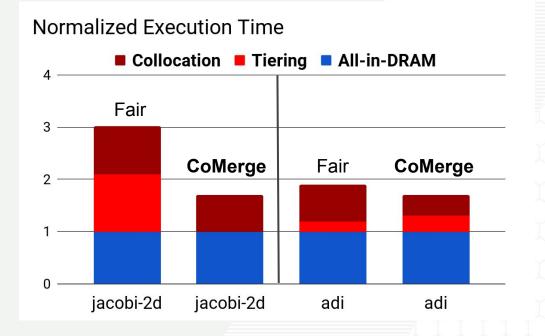
## **Sharing DRAM**

Sorting objects using co-benefit metric.









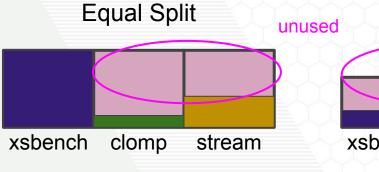


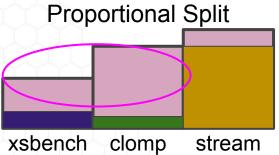
### Summary

### More detailed analysis in the paper



Partitioning & existing solutions





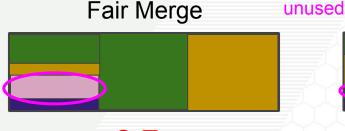
7x

slowdown

6x

CoMerge

Sharing & co-benefit metric



2.7x

slowdown

2.6x

#### Co-Benefit metric allows CoMerge to achieve:

- Lower runtime across all collocated applications.
- Higher DRAM utilization.